

A Formal Foundation for Trace-Based JIT Compilation

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Problem: design formal framework for trace-based JIT compilation

Trace-based JIT Compilation

```
(define (fac n)
  (if (< n 2)
    1
    (* n
        (fac (- n 1)))))
```

Guo and Palsberg (2011)
Disegna et al. (2014)

Tracing tied to specific execution model

```
(label 'fac-loop)
(literal-value 2)
(save-val)
(lookup-var 'n)
(save-val)
(lookup-var '<)
(apply-native 2)
(guard-false)
(literal-value 1)
(save-val)
(lookup-var 'n)
(save-val)
(lookup-var '-')
(apply-native 2)
...
(goto 'fac-loop)
```

Language semantics
(interpreter)

From
Local optimisations
on traces

Research goal

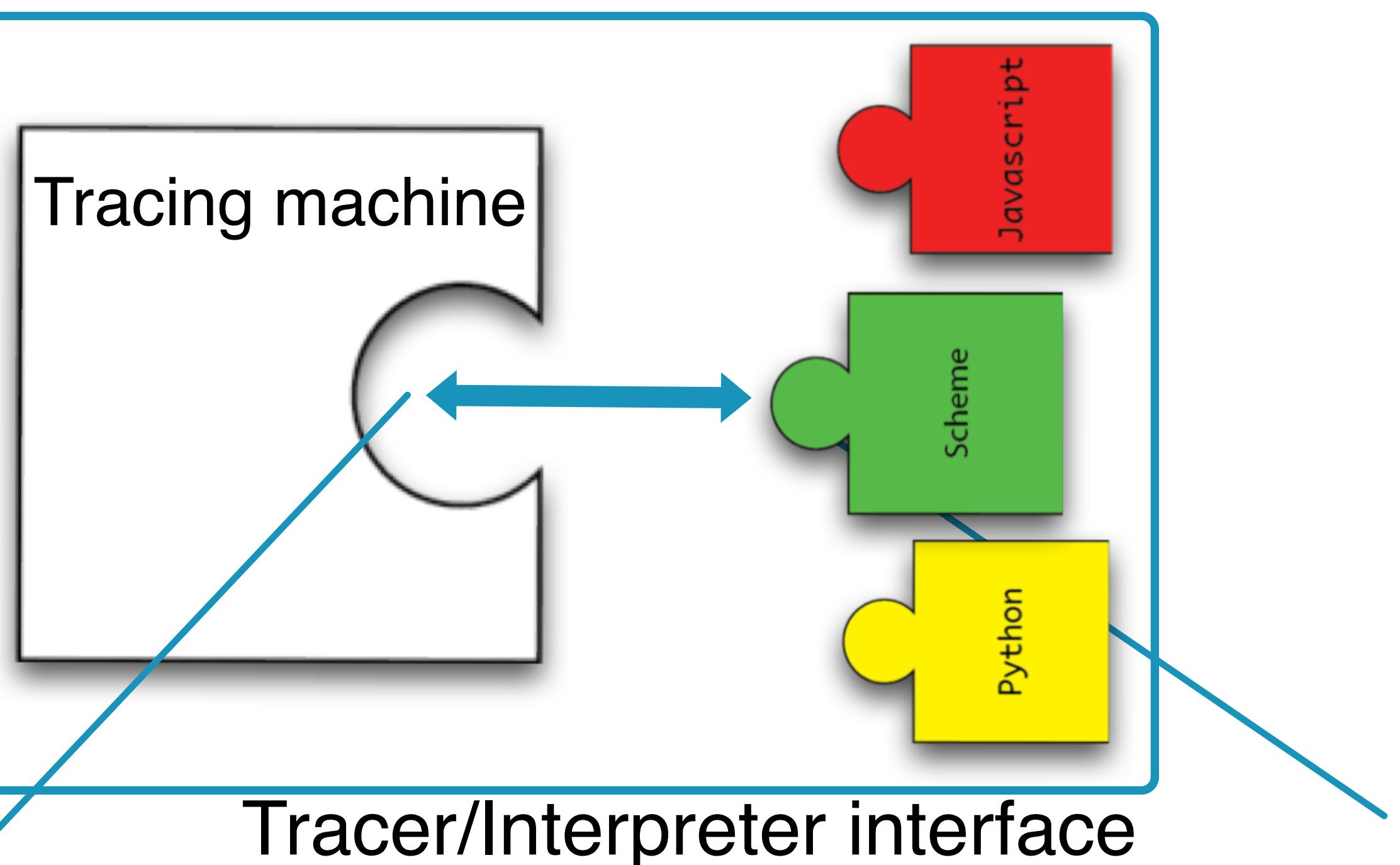
Traced language semantics
(TJIT Compiler)

Common framework
for trace and
non-trace execution

Recording/Executing
traces
Handling guard failures
Dynamic analysis hook

Towards
Global optimisations
between traces

Approach: separate tracer and interpreter by creating an explicit interface

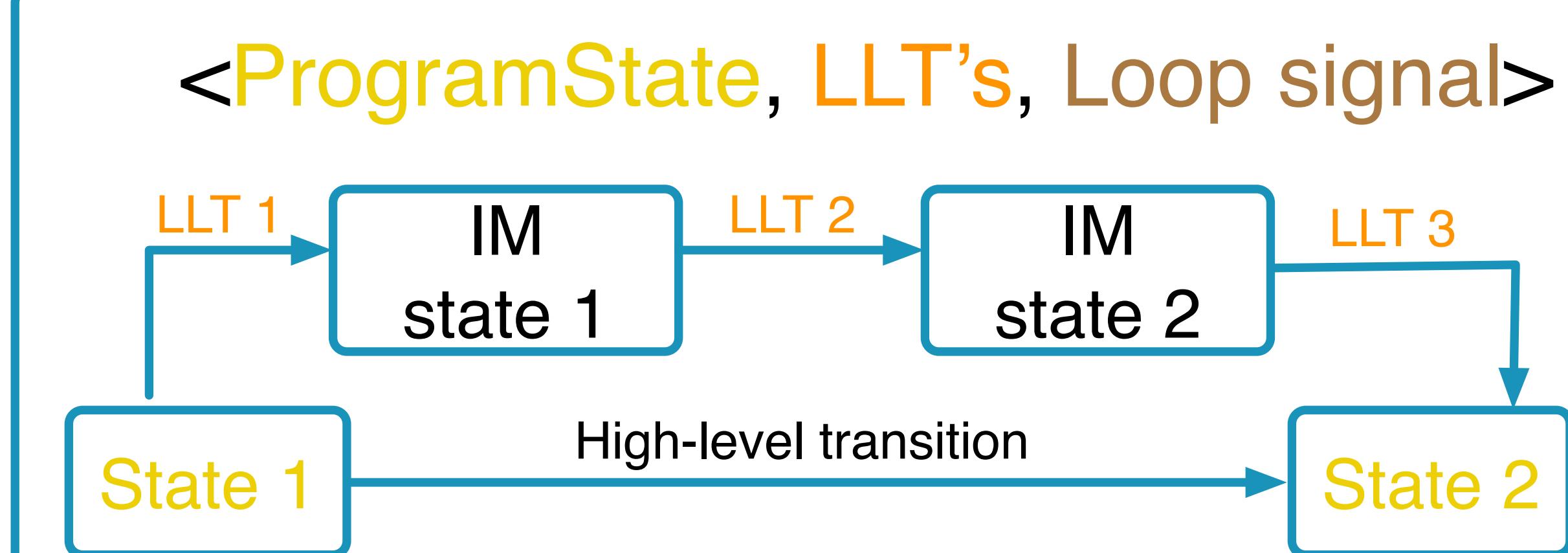


Interpreter as state machine

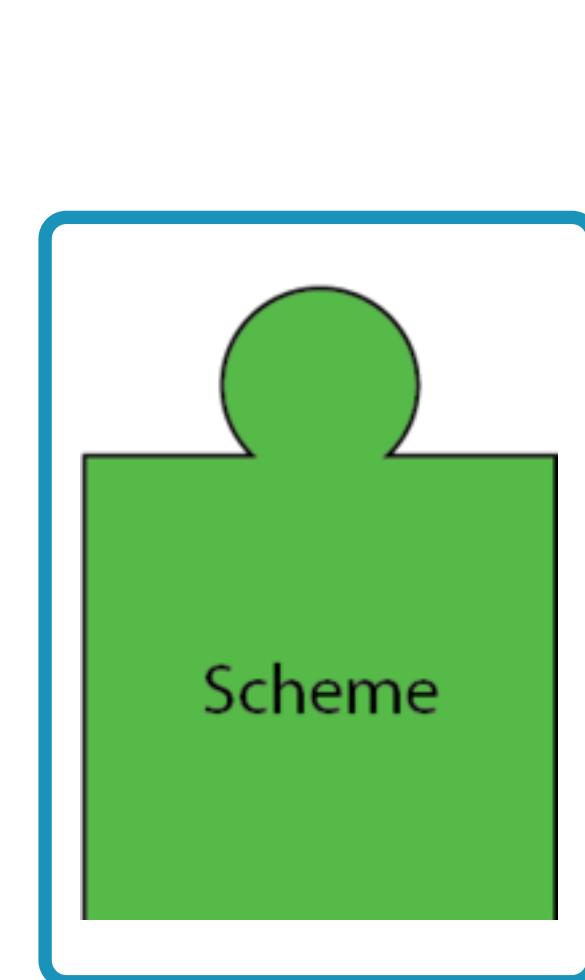
optimize: Trace → Trace

restart: RestartPoint X ProgramState
→ ProgramState

step: ProgramState → InterpreterReturn



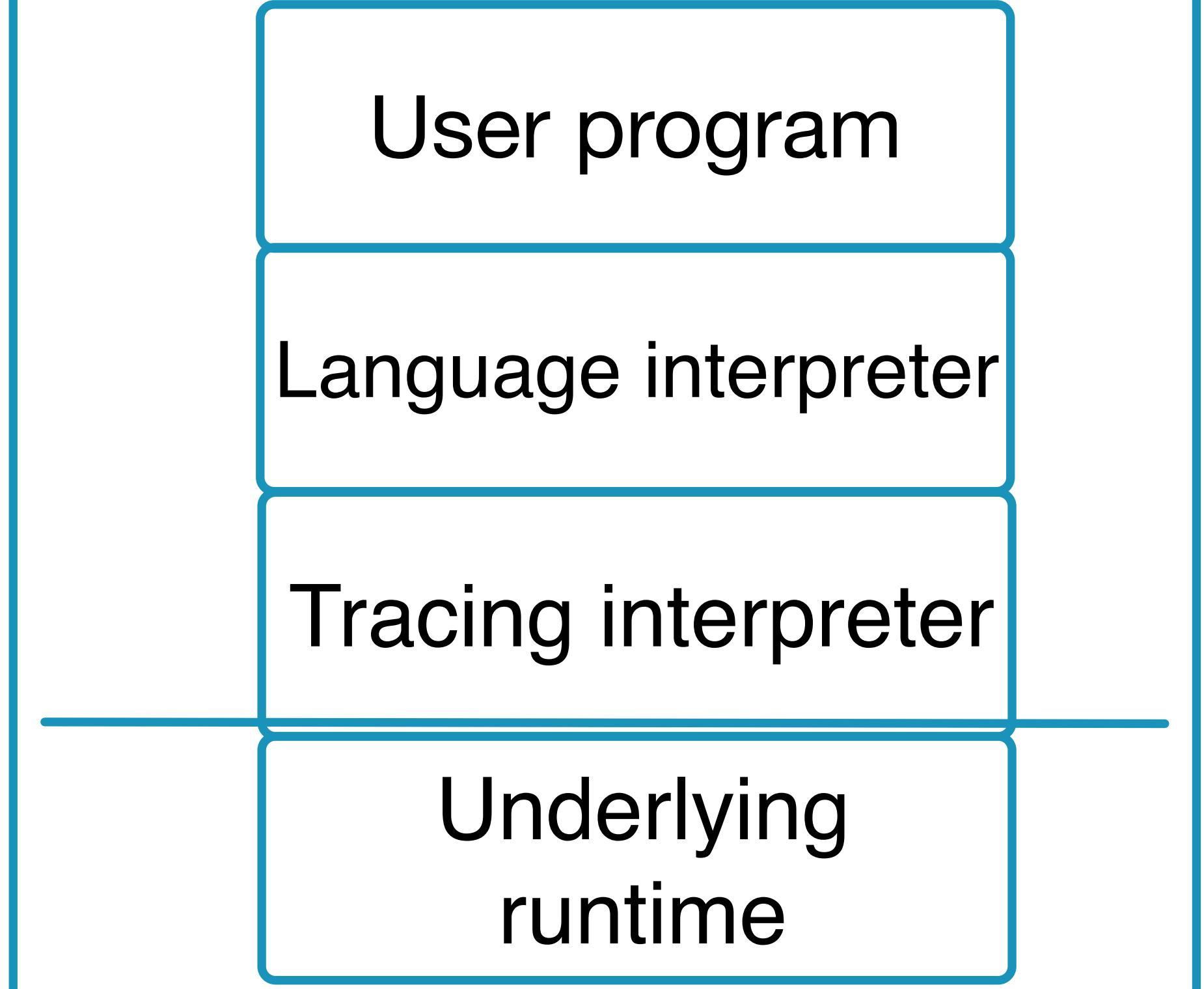
Experiment: apply transformation on existing interpreters



```
(match program-state
  ...
  ((ifk condition consequent alternative) p σ κ)
  (if condition
      (interpreter-return
        program-state
        (ev consequent)
        #f
        (restore-env)
        (pop-continuation)
        (guard-true alternative)))
  (interpreter-return
    program-state
    (ev alternative)
    #f
    (restore-env)
    (pop-continuation)
    (guard-false consequent))))
  ...
)
```

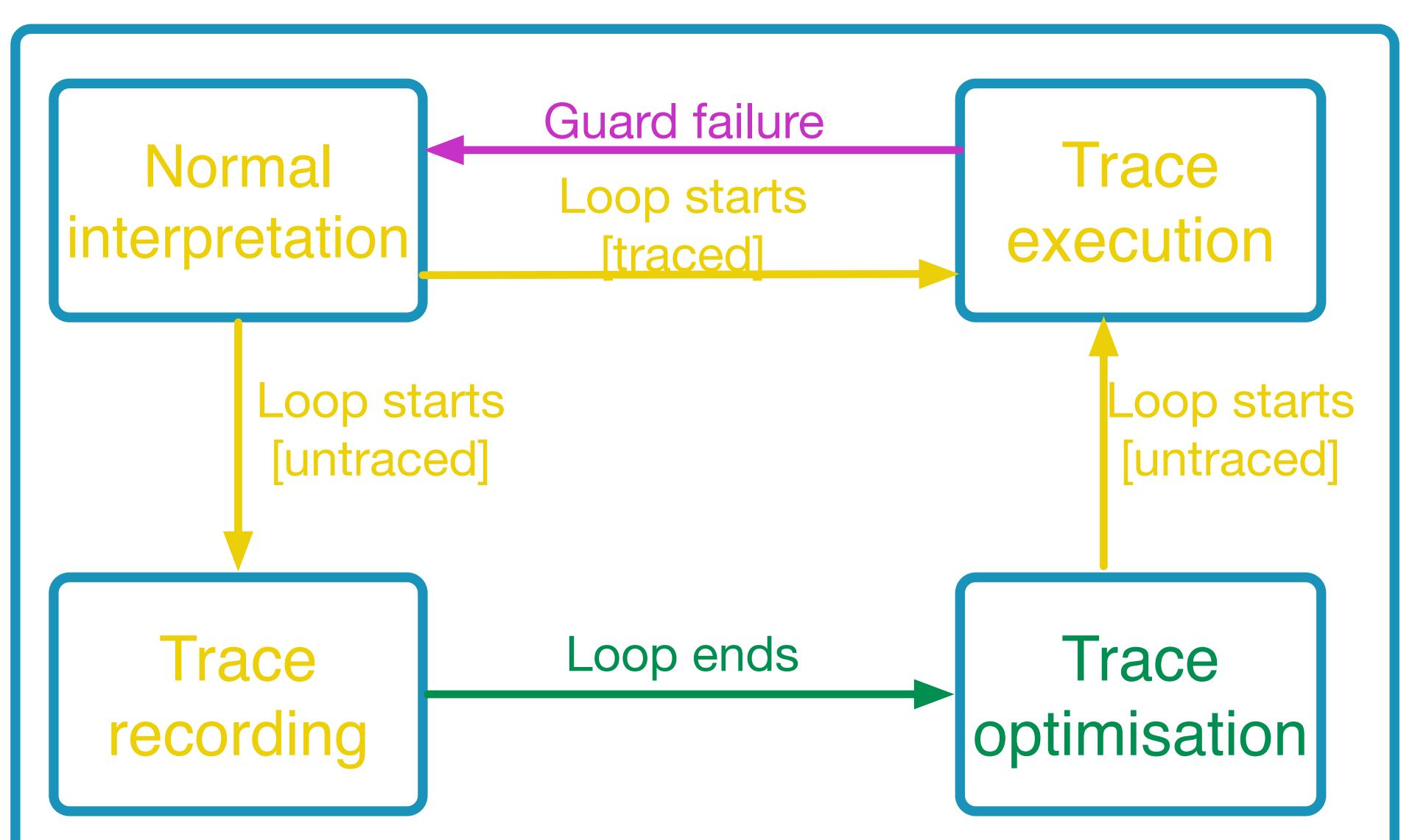
Minimal changes to implementation required

Meta-tracing compiler



Implementation: formalism + implementation

State machine



Formal semantics

```
ts(TE, tc, Σ, tn(lbl, τ)) →
  ts(TE, tc, Σ', tn(lbl, τ))
  if τ(Σ) = traceStep(Σ')
ts(TE, tc, Σ, tn(lbl, τ)) →
  ts(NI, tc, Σ', False)
  if τ(Σ) = guardFailed(rp)
  and where Σ' = restart(rp, Σ)
ts(TE, tc, Σ, tn(lbl, '() )) →
  ts(TE, tc, Σ, tn(lbl, τ))
  and where τ is the trace that has already been recorded for lbl
```

Implementation

```
define do-trace-executing-step
  let* ((trace-node (evaluator-state-trace-executing
                     evaluator-state))
        (trace (trace-node-trace trace-node))
        (label (trace-key-label (trace-node-trace-key
                                 trace-node))))
  if (null? trace)
    let* ((old-trace-node (evaluator-state-trace-executing
                           evaluator-state))
          (old-trace-key (trace-node-trace-key old-trace-node))
          (trace-key (make-label-trace-key
                      (trace-key-label old-trace-key)
                      (trace-key-debug-info old-trace-key)))
          (new-trace-node (get-trace-tracer-context trace-key)))
    (evaluator-state-copy evaluator-state
      (trace-executing new-trace-node)))
  let* ((instruction (car trace))
        (program-state
          (evaluator-state-program-state evaluator-state)))
    (handle-response-executing (instruction program-state))))
```