

Glitch: a programming model for Live Programming

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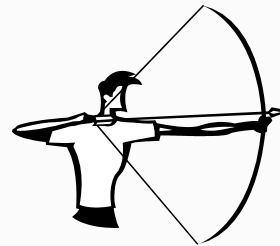
Live Feedback: Archer Analogy

[Hancock, 2003]

Archer:

aim → shoot → see → repeat ↻

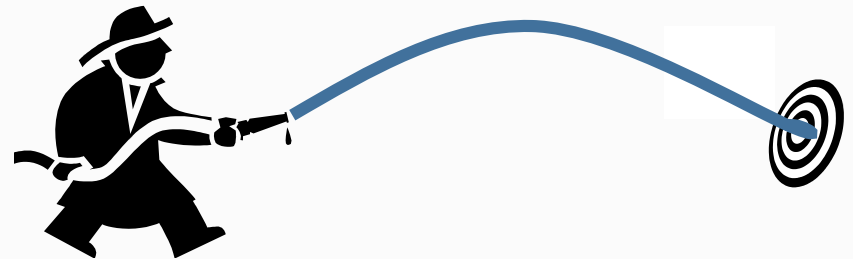
edit → run → debug → repeat ↻



Hose:

aim ↔ see

edit ↔ debug





Yin

Yang

Can we make this **real**?
(beyond small demos)

Challenges

Infrastructure (lots of change!)

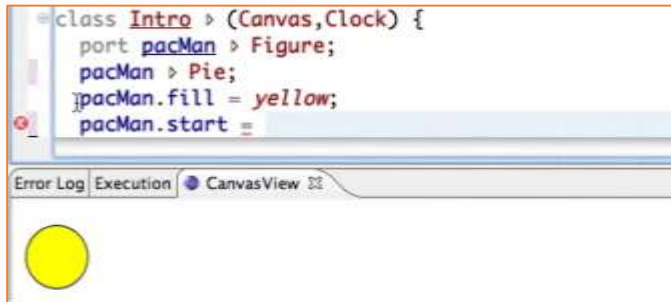
Incremental compiler

Editor

Execution engine

Living it up with a **Live Programming** Language

[McDermid, 2007]



```
class Intro > (Canvas, Clock) {  
  port pacMan > Figure;  
  pacMan > Pie;  
  }pacMan.fill = yellow;  
  pacMan.start =
```

The screenshot shows a code editor with the above code. Below the code is a control panel with tabs for 'Error Log', 'Execution', and 'CanvasView'. The 'CanvasView' tab is active, displaying a yellow circle on a white background.

Superglue (my thesis)

FRP-like language with signals

Handling code changes “fit” into its programming model

Cool **toy** demos

What about **real** code!?

All infrastructure for this demo was written with something else!

Living it up with a **Live Programming** Language

[McDirmid, 2007]

Incremental Compilation Framework

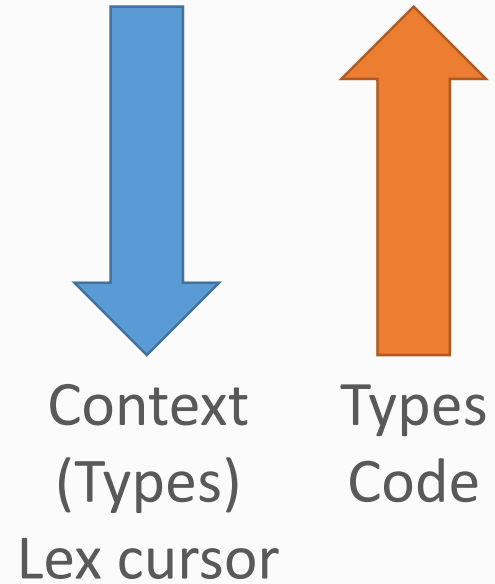
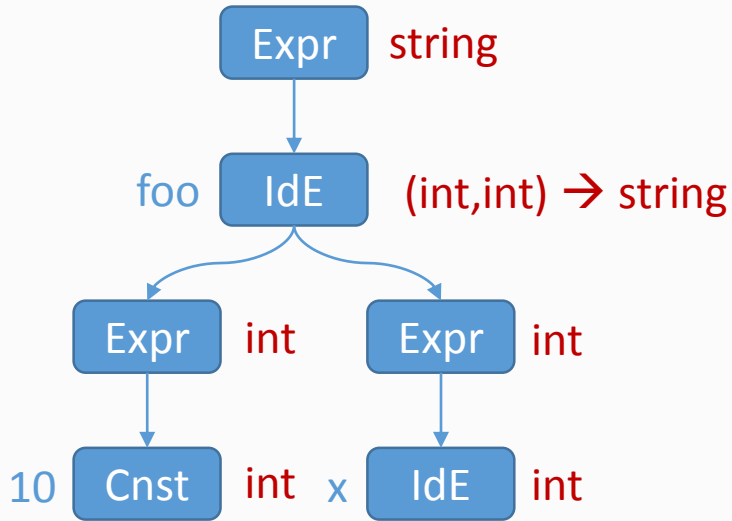
Damage/repair of memoized tree nodes

Originally for **Scala** in **Scala**



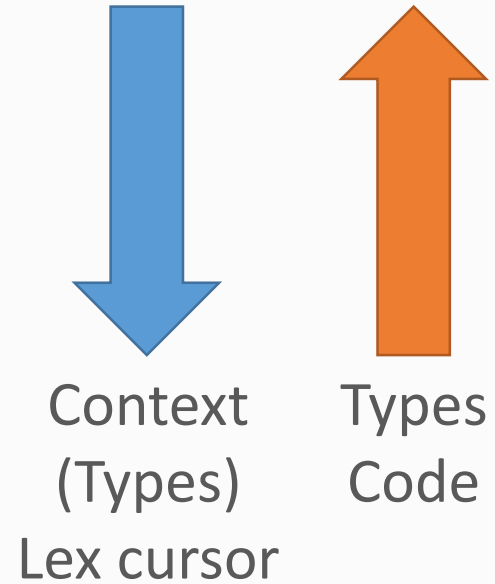
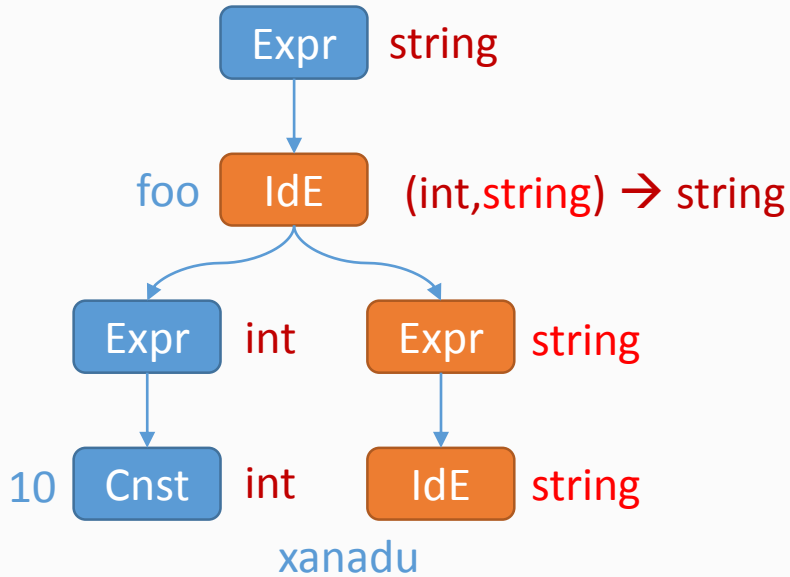
Incremental Compilation

foo(10, x)



Incremental Compilation

foo(10, xanadu)



Incremental Compilation

Type checkers need **symbol tables**

Indirection, key names depend on input, non-local

Go beyond plain **memoization**

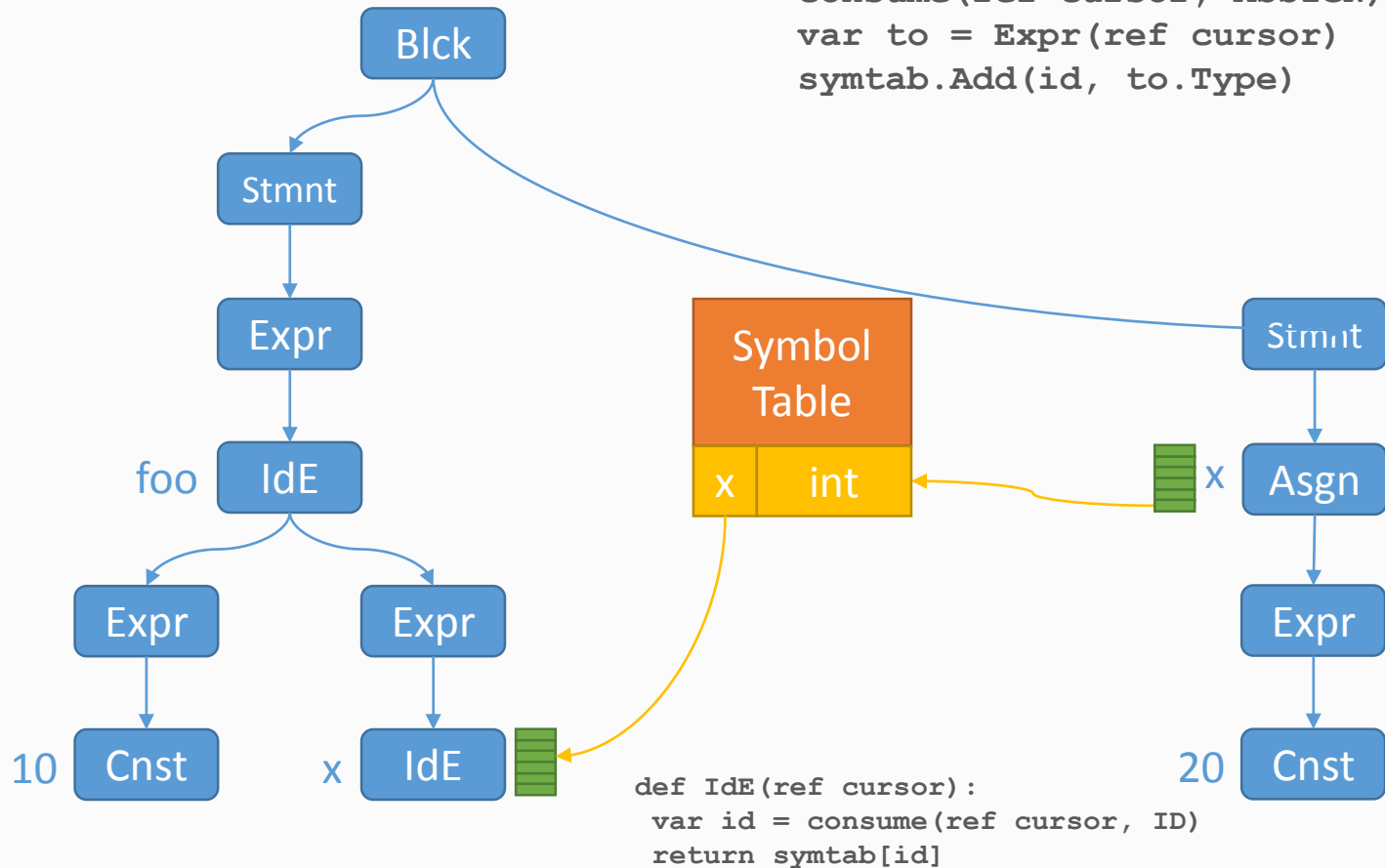
Trace non-local dependencies

Log, undo symbol table entries per tree

Incremental Compilation

foo(10, x)
x = 20

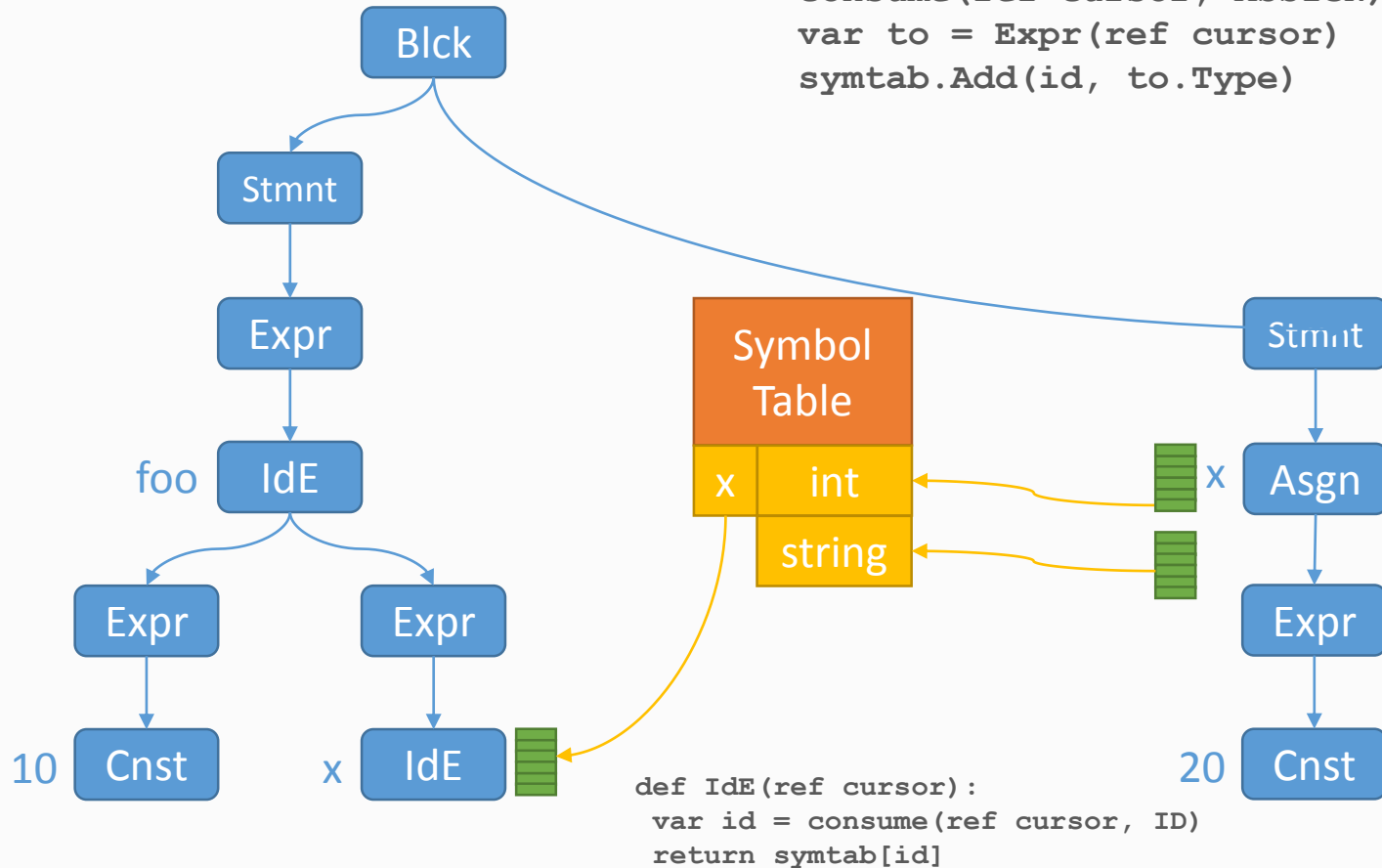
```
def Assign(ref cursor):  
    var id = consume(ref cursor, ID)  
    consume(ref cursor, ASSIGN)  
    var to = Expr(ref cursor)  
    symtab.Add(id, to.Type)
```



Incremental Compilation

foo(10, x)
x = "hello"

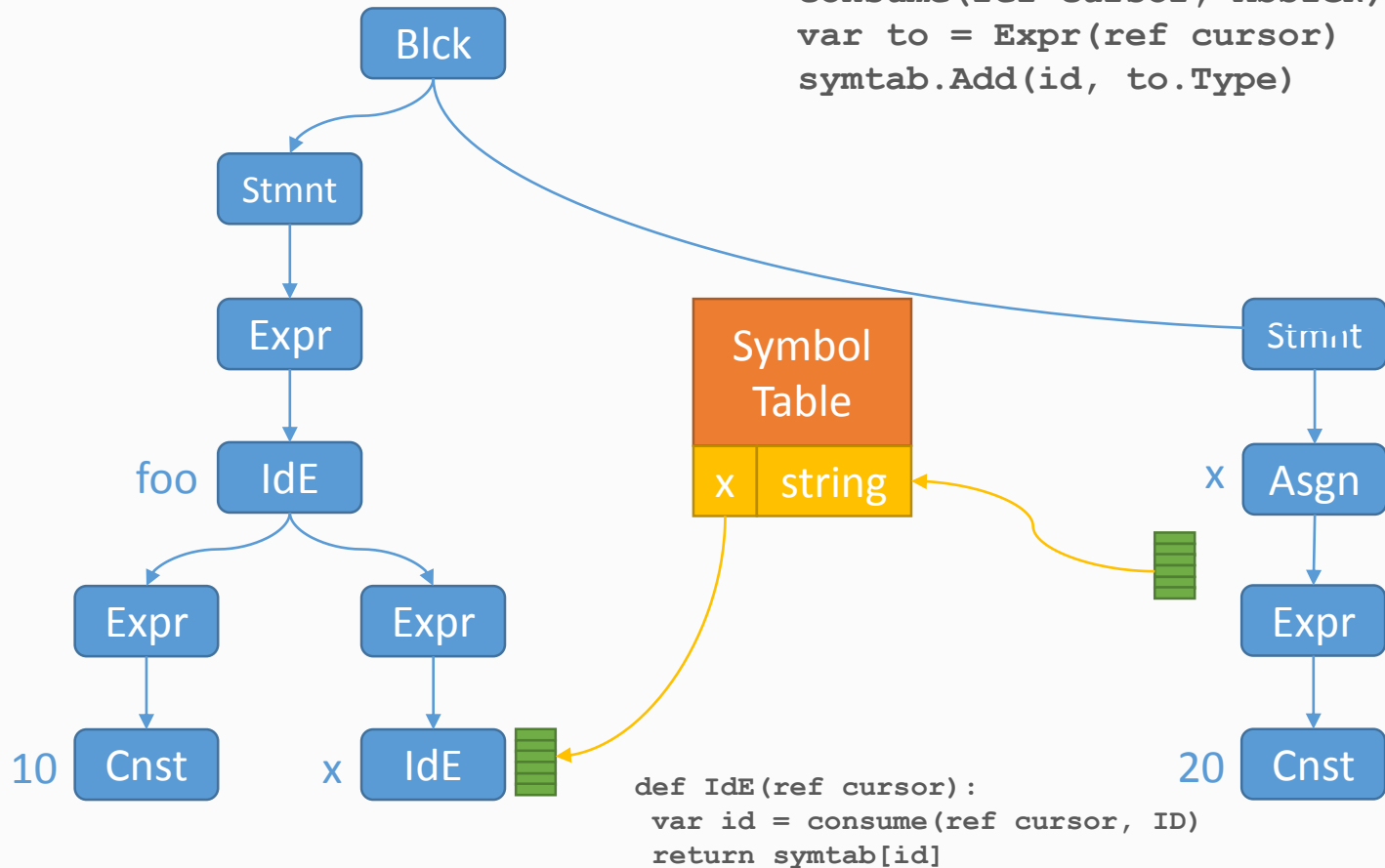
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Incremental Compilation

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Incremental **Compilation**

Able to weave this into scalac

Could mostly handle the way scalac was implemented

Actually this was a requirement...

No need to **lose** imperative programming

But then scalac also was not very imperative...

Adapted quickly for SuperGlue



Usable Live Programming

[McDirmid, 2013]

Live programming in 2013 is cool again!

Time to dust off my old tricks

This time...generalize

Infrastructure/language based on **same programming model**

“symbol table add” → any imperative operation?

Glitch as the simplest model
that could possibly work

Glitch

Divide program execution up into **nodes**

Trace dependencies

Log side-effecting operations (imperatives)

Re-execute **node** on dependency change

Reap after re-exec: undo **dead** imperatives

Compare old and new log; recursively undo dead nodes

Glitch

Imperatives must be undoable

Nodes can re-execute in **any order**:

Imperatives must be commutative w.r.t. order!

Supported imperatives:

Set Add, Aggregation (trivial)

Assignment (dynamic single assignment restriction)

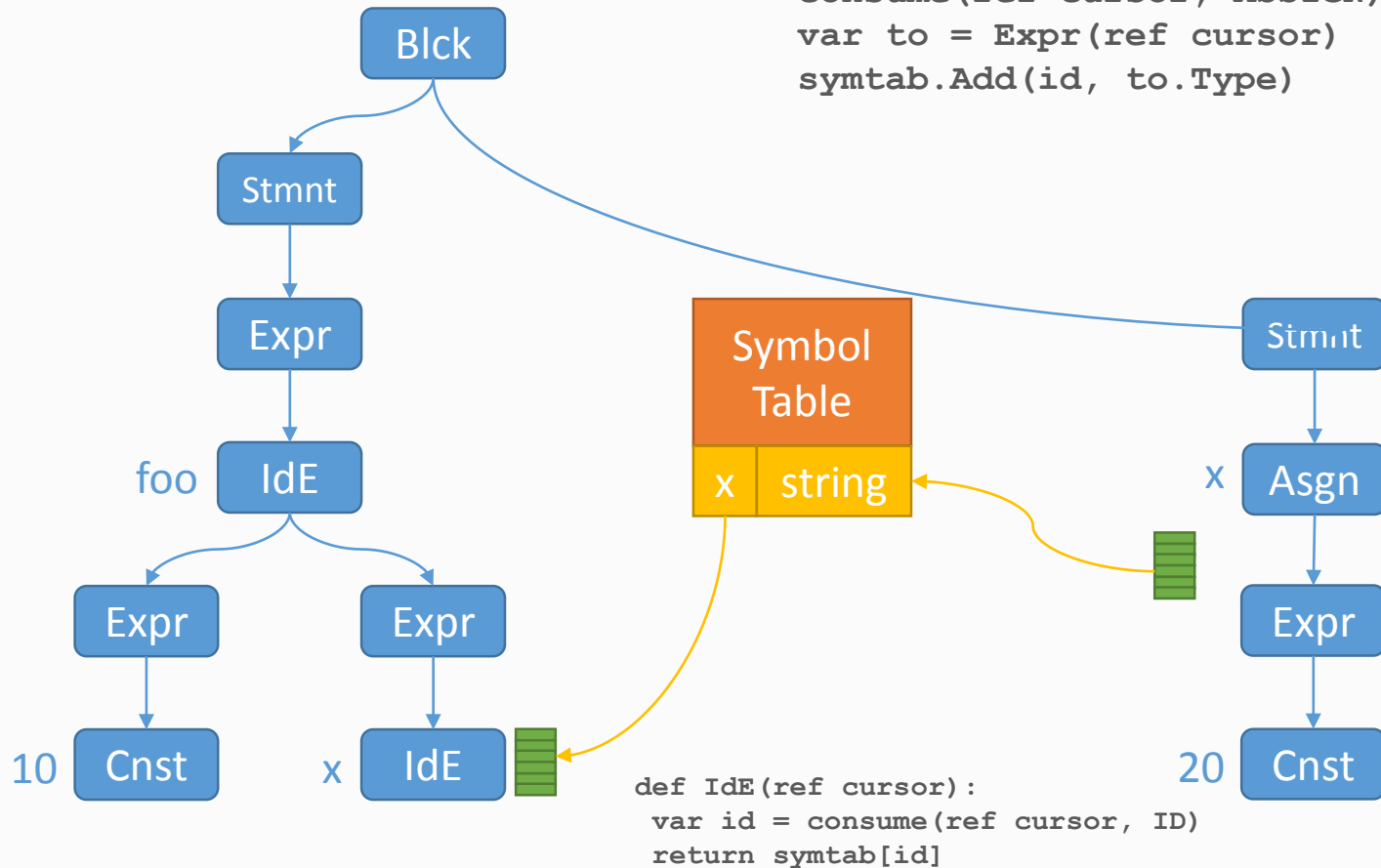
Dictionary Set (like assignment)

List.Append (provide ordered execution address)

Glitch

foo(10, x)
x = "hello"

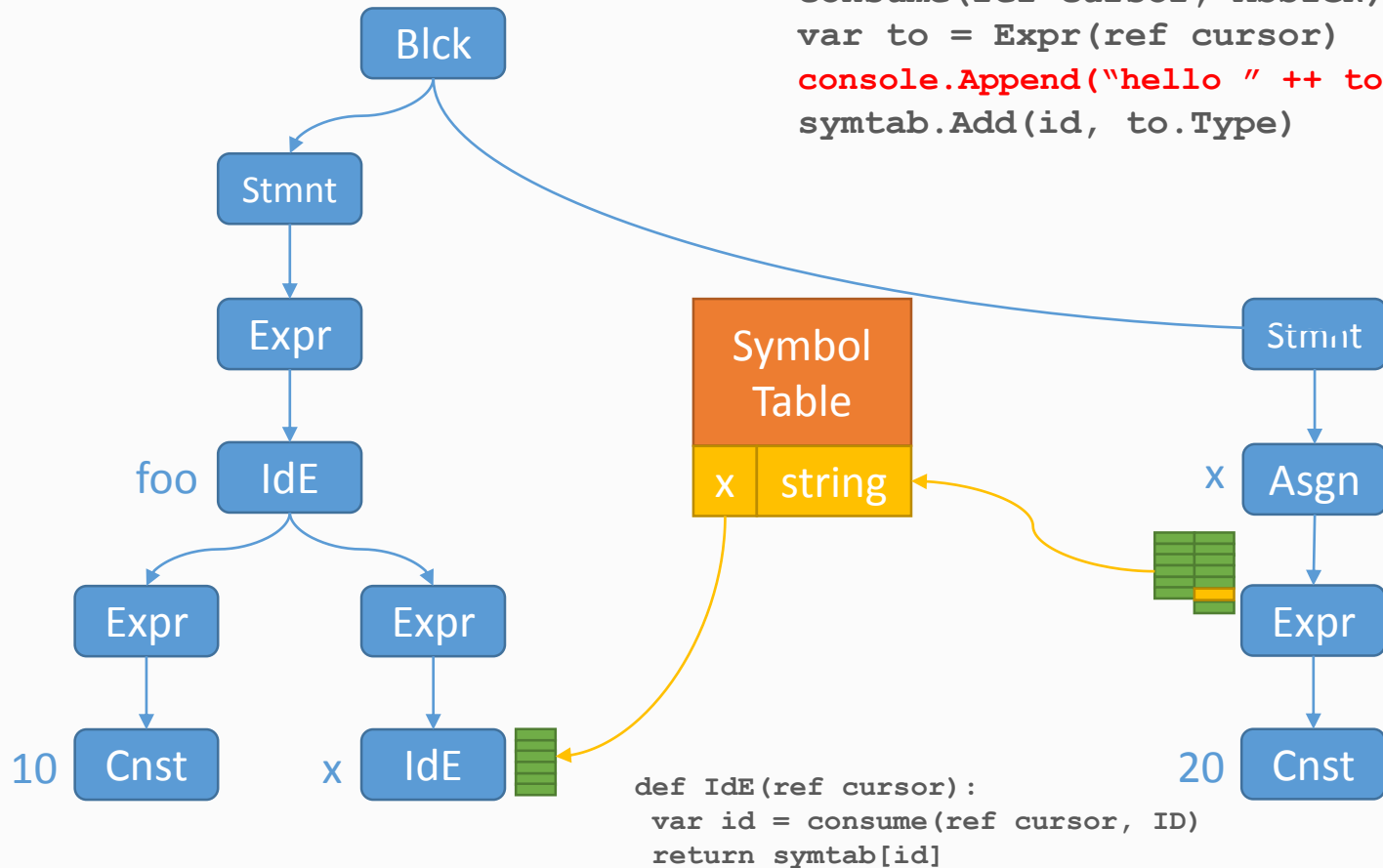
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    symtab.Add(id, to.Type)
```



Glitch

```
foo(10, x)  
x = "hello"
```

```
def Assign(ref cursor):  
    var id = consume(ref cursor, ID)  
    consume(ref cursor, ASSIGN)  
    var to = Expr(ref cursor)  
    console.Append("hello " ++ to.Tp)  
    symtab.Add(id, to.Type)
```



Why “Glitch”?

Consistency is **eventual**

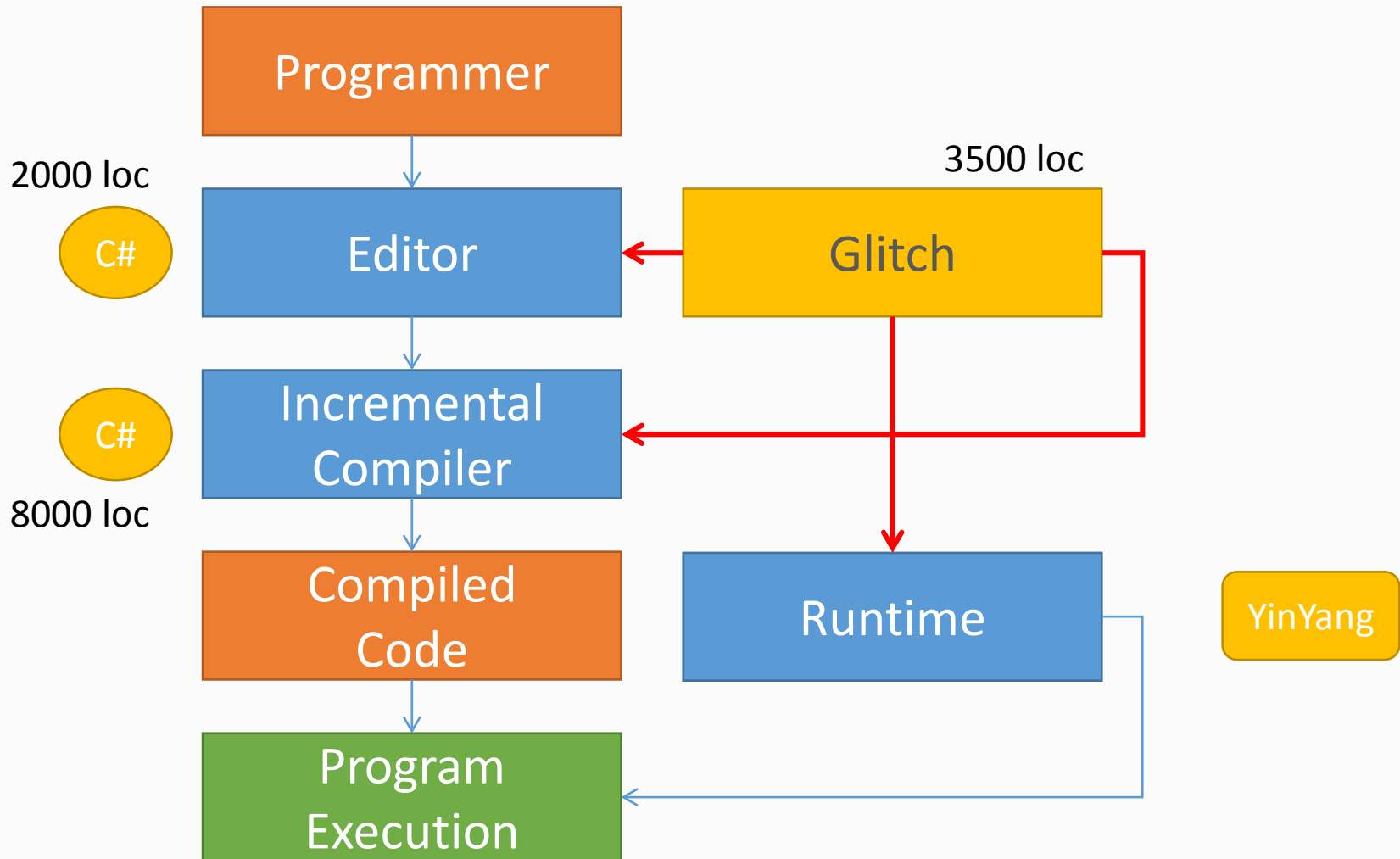
No attempt is made to find an “optimal” re-execution order

Sandbox during development, commit when nothing is damaged in production

No fancy analyses; just logging

Kind of **boring** from a technical perspective, but it works!

Glitch in the stack



Glitch payoffs

Expressive semi-imperative code

Automatic repair management

Simple implementation

Completely dynamic (also a cost...)

*But wait, there's **more!***

Iterative computing

No restriction on what writes can be seen

Parsing and type checking together in same pass

Optimistic speculative parallelism

Distributed computing

...

Asynchronous Execution

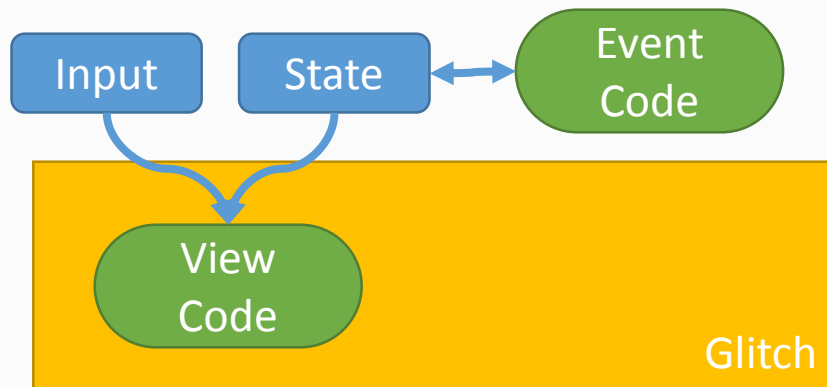
```
// synchronous - old/blocking/deterministic
var dataA = fileA.read()
process(dataA)
var dataB = fileB.read()
process(dataB)
```

```
// asynchronous - new/non-blocking/non-deterministic
fileA.readAsync(dataA => process(dataA))
fileB.readAsync(dataB => process(dataB))
```

```
// Glitch - retro/non-blocking/deterministic
var dataA = fileA.read()
process(dataA)
var dataB = fileB.read()
process(dataB)
```

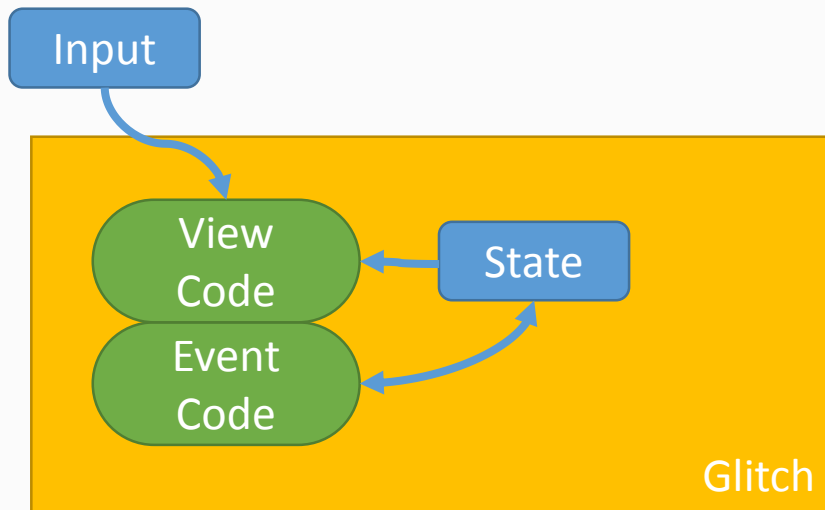
Time as major future work

Cannot abstract over time, events, and **interactivity**
...so no debugging of time-stepped computations (yet)



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Cannot abstract over time, events, and **interactivity**
...so no debugging of time-stepped computations (yet)



Drawbacks

Tracing logging have costs (performance)

Only semi-imperative (expressiveness)

Eventual consistency makes pulling
triggers harder

Not complete until time is included

Conclusion

Why not manage **change** like **garbage collection** manages memory?

