A Graph-Based Formal Semantics of Reactive Programming from First Principles

<u>**Bjarno Oeyen**</u>, Joeri De Koster, Wolfgang De Meuter Software Languages Lab, Vrije Universiteit Brussel, Belgium

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(defr (average x y z) (/ (+ x y z) 3))



- Syntactically looks like Scheme code
 - Makes graphs
- Reactive semantics
 - Instead of *applying* functions, reactors are *deployed*.
 - Instances of reactors are called **deployments**

Push-based propagation

- Evaluation of RP programs in **turns**.
- Haai-specific characteristics
 - Everything is a reactor (even + and /)
 - Higher-Order Reactors

Higher-Order Sounds like... Haai Is Dutch for... Shark



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Reactive Programming in Haai

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Push-based propagation

Evaluation of RP programs in **turns**. •

Haai-specific characteristics

- Everything is a reactor (even + and /)
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Reactive semantics

• Makes graphs

- Instead of *applying* functions, reactors are *deployed*.
- Instances of reactors are called deployments







Reactive Programming in a nutshell

- Automatic recomputation of program state
 - By **declaring** constraints/dependencies between *signals*
 - No callback hell to keep data dependencies updated
- Sources bound to data producers
 - E.g., user input, sensors...
- Sinks bound to data consumers
 - E.g., actuators, user interface...
- Other "RP" languages: REScala, Frappé, FrTime, Elm, ReactiveX, Akka Streams...



Motivation

Formalisation

Lessons learned

Two implementation styles of RP languages

Function-Based

Functions & Function Composition Examples: Fran, Yampa, SFRP, Dunai...

• Usually implemented in Haskell...

Well-studied (∃ many formalisations of function-based RP)

So why bother formalising graphbased RP?

Graph-Based

Graphs & Graph Composition Examples: FrTime, REScala, Frappé, EmFRP... Not as well studied as function-based RP

Memory Usage of a Yampa program



- Haskell RP languages are, in general, unsuitable for embedded devices [*]
 - E.g., IoT, CPS, Real-Time Systems
- Risk of space leaks
- Need for a garbage collector
- •
- Which is usually also the case of their formalisations.

Motivation

Formalisation

Lessons learned

Karcharias









(defr (average x y z) (/ (+ x y z) 3))

(average sensor0 sensor1 sensor2)

Construct a *cascade* of reactor deployments (instances) to fully connect the average reactor.

(defr (average x y z) (/ (+ x y z) 3))

(average sensor0 sensor1 sensor2)

average ($\iota_{d,avg}$)		
Name	Signal	
in _{1,0}		
in _{2,0}		
in _{3,0}		
s ₀		
out _{1,0}		



(defr (average x y z)

(/ (+ x y z) 3))



+ (<i>ι_{d,+}</i>)	
Name	Signal
in _{1,0}	
in _{2,0}	
in _{3,0}	
out _{1,0}	

(defr (average x y z)

(/ (+ x y z) 3))



(defr (average x y z) (/ (+ x y z) 3))















Wiring (W-Rules)

- Connect signals between deployments
- Signal environr by-step
 - Set of node

(W-REF)

$$w \in W ::= \mathcal{W} \langle \iota_d, N, \Sigma \rangle$$

Deployment identifier \checkmark
Node set \checkmark
Signal environment \checkmark

(W-CONGRUENCE) $w \rightarrow_w w'$ $\mathcal{K}\langle E, I_d, \{w\} \cup W, S, D \rangle \to_k \mathcal{K}\langle E, I_d, \{w'\} \cup W, S, D \rangle$

Global Rules (\rightarrow_k) :

Local Wiring Rules (\rightarrow_{W}) :

 $\mathcal{W}\langle \iota_d, \{(\bar{i}_l \S[x] \S \bar{i}_r, nt, \bar{o})\} \cup N, \Sigma \rangle$ $\rightarrow_{\mathcal{W}} \mathcal{W}(\iota_d, \{(\bar{i}_l \, \{ [\Sigma(x)] \, \{ \bar{i}_r, nt, \bar{o} \} \} \cup N, \Sigma \})$

$$w_{avg} = \mathcal{W} \langle \iota_{d,avg}, \\ \{([+, in_{1,0}, in_{2,0}, in_{3,0}], \mathcal{DEPLOY}, [x]) \\ ([/, x, 3], \mathcal{DEPLOY}, [out_{1,0}]) \}, \\ \{in_{1,0} \mapsto S_{\mathcal{REF}} \langle \dots \rangle, \dots \\ time \mapsto S_{\mathcal{GLB}} \langle time \rangle, \dots \} \rangle$$

$$w'_{avg} = \mathcal{W} \langle \iota_{d,avg}, \\ \{([/, x, 3], \mathcal{DEPLOY}, [out_{1,0}]) \}, \\ \{x \mapsto S_{\mathcal{REF}} \langle S_{\mathcal{REF}} \langle \iota_{d,avg}, d_1 \rangle, out_{1,0} \rangle, \\ in_{1,0} \mapsto S_{\mathcal{REF}} \langle \dots \rangle, \dots \\ time \mapsto S_{\mathcal{GLB}} \langle time \rangle, \dots \} \rangle$$

$$w''_{avg} = \mathcal{W} \langle \iota_{d,avg}, \\ \{\}, \\ \{out_{1,0} \mapsto S_{\mathcal{REF}} \langle S_{\mathcal{REF}} \langle \iota_{d,avg}, d_2 \rangle, out_{1,0} \rangle, \\ x \mapsto S_{\mathcal{REF}} \langle S_{\mathcal{REF}} \langle \iota_{d,avg}, d_1 \rangle, out_{1,0} \rangle, \\ in_{1,0} \mapsto S_{\mathcal{REF}} \langle S_{\mathcal{REF}} \langle \iota_{d,avg}, d_1 \rangle, out_{1,0} \rangle, \\ in_{1,0} \mapsto S_{\mathcal{REF}} \langle \dots \rangle, \dots \\ time \mapsto S_{\mathcal{GLB}} \langle time \rangle, \dots \} \rangle$$





















Evaluation Contexts:

Propagation (S-Rules)

- Reduce signal's into their current values.
- S-rules define the semantics of computing a **snapshot**.

$$s \in S ::= S \langle \iota_d, \Sigma \rangle$$

Deployment identifier



It's not that easy ...

(defr (weird x y z) (def +or* ...) (/ (+or* x y z) 3))

Higher-Order Reactivity

Signal

 $\mathcal{S}_{GLB}(\text{sensor0})$

 $\mathcal{S}_{GLB}(\text{sensor1})$

 $\mathcal{S}_{\mathcal{GLB}}(\text{sensor2})$

 $\mathcal{S}_{\mathcal{REF}}\langle \mathcal{S}_{\mathcal{REF}}\langle \iota_{d,avg}, d_0 \rangle, out_{1,0} \rangle$

•••

Wiring decis

at-propagation

at wirin

average ($\iota_{d,avg}$)

Name

in_{1,0}

in_{2,0}

in_{3,0}

S₀

•••

	+ (<i>ι_{d,+}</i>)	
	Name	Value
	in _{1,0}	$\mathcal{S}_{\textit{REF}}\langle \iota_{d,avg}, in_{1,0} angle$
	In _{2,0}	$\mathcal{S}_{\mathrm{REF}}\langle \iota_{d,avg}, in_{2,0} angle$
	in _{3,0}	$\mathcal{S}_{\mathrm{REF}}\langle \iota_{d,avg}, in_{3,0} angle$
on is made		
g time.	out _{1,0}	$\mathcal{S}_{\mathcal{DEP}}\langle\iota_{b,+}, [+],\ldots angle$
	* (<i>ι_{d,*}</i>)	
	* (ι _{d,*}) Name	Value
	* (<i>L</i> _{<i>d</i>,*}) <i>Name</i> in _{1,0}	Value $\mathcal{S}_{\mathcal{REF}}\langle \iota_{d,avg},in_{1,0} angle$
	* (<i>ι_{d,*}</i>) <i>Name</i> in _{1,0} In _{2,0}	Value $S_{REF} \langle \iota_{d,avg}, in_{1,0} \rangle$ $S_{REF} \langle \iota_{d,avg}, in_{2,0} \rangle$
	* $(\iota_{d,*})$ <i>Name</i> $in_{1,0}$ $ln_{2,0}$ $in_{3,0}$	Value $S_{REF} \langle \iota_{d,avg}, in_{1,0} \rangle$ $S_{REF} \langle \iota_{d,avg}, in_{2,0} \rangle$ $S_{REF} \langle \iota_{d,avg}, in_{3,0} \rangle$
	* $(\iota_{d,*})$ <i>Name</i> $in_{1,0}$ $ln_{2,0}$ $in_{3,0}$	Value $S_{REF} \langle \iota_{d,avg}, in_{1,0} \rangle$ $S_{REF} \langle \iota_{d,avg}, in_{2,0} \rangle$ $S_{REF} \langle \iota_{d,avg}, in_{3,0} \rangle$

(defr (weird x y z) (def +or* ...) (/ (+or* x y z) 3))

+ (*ι*_{d,+})

				Name	Value
				in _{1,0}	$\mathcal{S}_{\mathcal{REF}}\langle \iota_{d,avg}, in_{1,0} angle$
	average ($\iota_{d,avg}$)		In _{2,0}	$\mathcal{S}_{\mathrm{REF}}\langle \iota_{d,avg}, in_{2,0} \rangle$
	Name	Signal		in _{3,0}	$\mathcal{S}_{\mathcal{REF}}\langle \iota_{d,avg}, in_{3,0} \rangle$
	in _{1,0}	$\mathcal{S}_{\mathcal{GLB}}(\mathrm{sensor}0)$			
	in _{2,0}	$\mathcal{S}_{\mathcal{GLB}}(ext{sensor1})$		out _{1,0}	$\mathcal{S}_{\mathcal{DEP}}\langle\iota_{b,+}, [+],\ldots angle$
	in _{3,0}	$\mathcal{S}_{\mathcal{GLB}}(\mathrm{sensor2})$		*/. \	_
				* (<i>l_{d,*}</i>)	
Meta signal with the deployment identifier (ι_d)	d	$S_{arro}(\mu_{d}, S_{arro}(\mu_{d,arro} + -0r^{*}))$	2	Name	Value
	►u ₀			in _{1,0}	$\mathcal{S}_{\mathcal{REF}}\langle \iota_{d,avg}, in_{1,0} angle$
	S ₀	$\mathcal{S}_{\mathcal{REF}} \langle \mathcal{S}_{\mathcal{REF}} \langle t_{d,avg}, a_0 \rangle, Out_{1,0} \rangle$		In _{2,0}	$\mathcal{S}_{\textit{REF}} \langle \iota_{d,avg}, in_{2,0} \rangle$
				in _{3,0}	$\mathcal{S}_{\mathcal{REF}}\langle \iota_{d,avg}, in_{3,0} angle$
				out _{1,0}	

(defr (weird x y z) (def +or* ...) (/ (+or* x y z) 3))

+ (*ι*_{d,+})

				Name	Value	
				in _{1,0}	$\mathcal{S}_{REF}\langle \iota_{d,avg}, in_{1,0} angle$	
	average ($\iota_{d,avg}$,)		In _{2,0}	$\mathcal{S}_{\textit{REF}}\langle \iota_{d,avg}, in_{2,0} angle$	
	Name	Signal		in _{3,0}	$\mathcal{S}_{\textit{REF}}\langle \iota_{d,avg}, in_{3,0} \rangle$	
	in _{1,0}	<mark>د 10 s</mark> ensor0				
	in _{2,0}	<mark>کی</mark> sensor1)		out _{1,0}	$\mathcal{S}_{\mathcal{DEP}}\langle\iota_{b,+}, {}^{f+l},\ldots angle$	
	in _{3,0}	<mark>ζ₃₀sensor2</mark>)		* / . \	_	
				* (<i>l_{d,*}</i>)		
Meta signal with the deployment identifier (ι_d)	d	$S_{\sigma\pi\sigma}(l_{h,d}, S_{\sigma\pi\pi}(l_{d,ano}, +-0r^{*}), \ldots)$?	Name	Value	
	u ₀			in _{1,0}	$\mathcal{S}_{\mathcal{REF}}\langle \iota_{d,avg}, in_{1,0} angle$	
	S ₀	$\mathcal{S}_{\mathcal{R}EF} \langle \mathcal{S}_{\mathcal{R}EF} \langle \iota_{d,avg}, u_0 \rangle, \mathcal{O}u\iota_{1,0} \rangle$		In _{2,0}	$\mathcal{S}_{\mathcal{REF}}\langle \iota_{d,avg}, in_{2,0} angle$	
				in _{3,0}	$\mathcal{S}_{\mathcal{REF}}\langle \iota_{d,avg}, in_{3,0} angle$	
				out _{1,0}		

(defr (weird x y z) (def +or* ...) (/ (+or* x y z) 3))



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(defr (weird x y z) (def +or* ...) (/ (+or* x y z) 3))

Higher-Order Reactivity

average ($\iota_{d,avg}$)		
Name	Signal	
in _{1,0}	د 10 sensor0	
in _{2,0}	ک sensor1)	
in _{3,0}	ع sensor2	
d ₀	$\mathcal{S}_{\mathcal{D}\!E\!\mathcal{I}}\langle \iota_{b,d}, \mathcal{S}_{\mathcal{R}\!E\!\mathcal{I}}\langle \iota_{d,avg}, + ext{-or-}^* angle, \ldots angle$	
s ₀	$\mathcal{S}_{\mathcal{REF}} \langle \mathcal{S}_{\mathcal{REF}} \langle \iota_{d,avg}, d_0 \rangle, out_{1,0} \rangle$	
d1	$\mathcal{S}_{\mathcal{D}\!E\!E\!P}\langle\iota_{b,d}',\mathcal{S}_{\mathcal{G}\!L\!B}\langle\prime angle,\ldots angle$	
out _{1,0}	$\mathcal{S}_{\mathcal{REF}}\langle\mathcal{S}_{\mathcal{REF}}\langle\iota_{d,avg},d_1 angle,out_{1,0} angle$	

what happens in the very first turn?



W-Rules and S-Rules can be interleaved!

• Dynamically create new deployments for new captures (reactors) on the operator signals.

(W-DEPLOY)

$$x, \iota_{b} \text{ fresh}$$

$$\frac{\Sigma' = \Sigma[x \mapsto S_{D \in \mathcal{P}} \langle \iota_{b}, \sigma, \overline{\sigma} \rangle][o_{i} \mapsto S_{\mathcal{R} \in \mathcal{F}} \langle S_{\mathcal{R} \in \mathcal{F}} \langle \iota_{d}, x \rangle, out_{i,0} \rangle | \forall i \in [1..|\overline{o}|]]}{W \langle \iota_{d}, \{([\sigma] \S \overline{\sigma}, D \in \mathcal{PLOY}, \overline{o})\} \cup N, \Sigma \rangle \to_{w} W \langle \iota_{d}, N, \Sigma' \rangle}$$

W-Rules and S-Rules can be interleaved!

• Dynamically create new deployments for new captures (reactors) on the operator signals.



Evaluation Contexts:
$$\mathcal{E}$$
::={ $x \mapsto \mathcal{E}_{\sigma}$ } $\cup \Sigma$ \mathcal{E}_{σ} :::= \Box $|$ $S_{\mathcal{REF}}(\mathcal{E}_{\sigma, x})$ $|$ $S_{\mathcal{REF}}(\mathcal{E}_{\sigma, x})$ $|$ $S_{\mathcal{REF}}(\mathcal{I}_b, \mathcal{E}_{\sigma}, \overline{\sigma})$ $\mathcal{S}_{\mathcal{DEP}}(\mathcal{I}_b, \mathcal{E}_{\sigma}, \overline{\sigma})$ $\mathcal{S}_{\mathcal{DEP}}(\mathcal{I}_b, \mathcal{P}, \overline{\sigma} \S[\mathcal{E}_{\sigma}] \S \overline{\sigma})$ (S-DEPLOY-NEW) $\Sigma = \mathcal{E}[S_{\mathcal{DEP}}(\mathcal{I}_b, \mathcal{C}(\mathcal{I}_c, N, \Sigma_c), \overline{\sigma})]$ $(\mathcal{I}_b, \mathcal{I}_c) \notin \operatorname{dom}(D)$ \mathcal{I}_d fresh $w = \mathcal{W}(\mathcal{I}_d, N, \Sigma_c[\operatorname{in}_{i,0} \mapsto \sigma_i \mid \forall i \in [1..|\overline{\sigma}|]])$ $\Sigma' = \mathcal{E}[\mathcal{I}_d]$ $D' = D[(\mathcal{I}_b, \mathcal{I}_c) \mapsto \mathcal{I}_d]$ $\mathcal{K}(\mathcal{E}, \mathcal{I}_d, W, \{\mathcal{S}(\mathcal{I}_d, \Sigma\}\} \cup \mathcal{S}, D)$ $\rightarrow_k \mathcal{K}(\mathcal{E}, \{\mathcal{I}_d\} \cup \mathcal{I}_d, \{w\} \cup W, \{\mathcal{S}(\mathcal{I}_d, \Sigma')\} \cup \mathcal{S}, D')$ (S-DEPLOY-EXISTING) $\Sigma = \mathcal{E}[S_{\mathcal{DEP}}(\mathcal{I}_b, \mathcal{C}(\mathcal{I}_c, N, \Sigma_c), \overline{\sigma})]$ $\mathcal{L}_d = D(\mathcal{I}_b, \mathcal{I}_c)$ $\Sigma' = \mathcal{E}[\mathcal{I}_d]$ $\mathcal{K}(\mathcal{E}, \mathcal{I}_d, W, \{\mathcal{S}(\mathcal{I}_d, \Sigma\}\} \cup \mathcal{S}, D)$ $\rightarrow_k \mathcal{K}(\mathcal{E}, \{\mathcal{I}_d'\} \cup \mathcal{I}_d, W, \{\mathcal{S}(\mathcal{I}_d, \Sigma'\}\} \cup \mathcal{S}, D)$

W-Rules and S-Rules can be interleaved!

Local Wiring Rules (\rightarrow_w) :

Dynamically create new deployments for new captures (reactors) on the operator signals.

- One reduction relation (\rightarrow_k)
 - Two helper relations (\rightarrow_w and \rightarrow_s)
 - Both local and global wiring and snapshot rules. •

(W-REF) $\mathcal{W}\langle \iota_d, \{(\overline{i}_l \S[x] \S \overline{i}_r, nt, \overline{o})\} \cup N, \Sigma \rangle$ $\rightarrow_{\mathcal{W}} \mathcal{W}(\iota_d, \{(\overline{i}_l \& [\Sigma(x)] \& \overline{i}_r, nt, \overline{o})\} \cup N, \Sigma\}$

(W-DEPLOY)

 x, ι_h fresh $\Sigma' = \Sigma[x \mapsto \mathcal{S}_{\mathcal{DEP}}(\iota_b, \sigma, \overline{\sigma})][o_i \mapsto \mathcal{S}_{\mathcal{REF}}(\mathcal{S}_{\mathcal{REF}}(\iota_d, x), out_{i,0}) | \forall i \in [1., |\overline{o}|]]$ $\mathcal{W}(\iota_d, \{([\sigma] \in \overline{\sigma}, \mathcal{DEPLOY}, \overline{o})\} \cup N, \Sigma) \to_{\mathcal{W}} \mathcal{W}(\iota_d, N, \Sigma')$

(W-RHO)

 ι_c fresh $\Sigma_c = \text{shift_io}(\Sigma)$ $c = \mathcal{C}(\iota_c, N_{inner}, \Sigma_c)$ $\mathcal{W}(\iota_d, \{(\overline{\sigma}, \mathcal{RHO}(N_{inner}), [\sigma])\} \cup N, \Sigma\} \to_{\mathcal{W}} \mathcal{W}(\iota_d, N, \Sigma[\sigma \mapsto c])$

Where shift_io(
$$\Sigma$$
) = { $x' \mapsto \sigma$ | $\forall x \mapsto \sigma \in \Sigma, x'$ =

$$\begin{cases}
in_{i,j+1} & x = in_{i,j} \\
out_{i,j+1} & x = out_{i,j} \\
x & otherwise
\end{cases}$$
Global Rules (\rightarrow_k):

(w-congruence)

 $w \rightarrow_w w'$ $\mathcal{K}\langle E, I_d, \{w\} \cup W, S, D \rangle \rightarrow_k \mathcal{K}\langle E, I_d, \{w'\} \cup W, S, D \rangle$

Evaluation Contexts:

```
ε
                                            \{x \mapsto \mathcal{E}_{\sigma}\} \cup \Sigma
\mathcal{E}_{\sigma}
                                             S_{REF}(\mathcal{E}_{\sigma}, x)
                                             S_{DEP}(\iota_b, \mathcal{E}_{\sigma}, \overline{\sigma})
                                            S_{\mathcal{DEP}}\langle \iota_b, p, \overline{\sigma} \S [\mathcal{E}_{\sigma}] \S \overline{\sigma} \rangle
```

Local Snapshot Rules (\rightarrow_s) :

(S-SELF-REF) (S-DEPLOY-PRIMITIVE) $\begin{array}{c} (\mathcal{E} \cup \mathcal{E} \cup \mathcal{E}) \\ \mathcal{S} \langle \iota_d, \mathcal{E} [\mathcal{S}_{\mathcal{D} E \mathcal{P}} \langle \iota_b, p, \overline{v} \rangle] \rangle \\ \rightarrow_s \mathcal{S} \langle \iota_d, \mathcal{E} [\delta_p(\overline{v})] \rangle \end{array}$ $\Sigma = \mathcal{E}[\mathcal{S}_{\mathcal{REF}}\langle \iota_d, x \rangle] \qquad v = \Sigma(x)$ $S\langle \iota_d, \Sigma \rangle \to_s S\langle \iota_d, \mathcal{E}[v] \rangle$

(S-TUPLE-REF) $S\langle \iota_d, \mathcal{E}[S_{\mathcal{REF}}\langle \overline{v}, out_{(0,i)}\rangle] \rangle$ $\rightarrow_{s} \mathcal{S}\langle \iota_{d}, \mathcal{E}[v_{i}] \rangle$

Global Rules (\rightarrow_k) :

(S-ACTIVATE) (S-CONGRUENCE) $\mathcal{W}(\iota_d, \emptyset, \Sigma) \in W$ $\iota_d \in I_d$ $\forall s \in S : s = S\langle i'_d, \Sigma' \rangle, i_d \neq i'_d$ $s \rightarrow s s'$ $\mathcal{K}\langle E, I_d, W, \{s\} \cup S, D \rangle$ $\mathcal{K}\langle E, I_d, W, S, D \rangle$ $\rightarrow_k \mathcal{K}\langle E, I_d, W, \{\mathcal{S}\langle \iota_d, \Sigma\rangle\} \cup S, D\rangle$ $\rightarrow_k \mathcal{K}\langle E, I_d, W, \{s'\} \cup S, D \rangle$

(S-REF)

 $\Sigma = \mathcal{E}[\mathcal{S}_{\mathcal{REF}}\langle u'_d, x \rangle] \qquad v = \Sigma'(x) \qquad \Sigma'' = \mathcal{E}[v]$ $\mathcal{K}\langle E, I_d, W, \{ \mathcal{S}\langle \iota_d, \Sigma \rangle, \mathcal{S}\langle \iota'_d, \Sigma' \rangle \} \cup S, D \rangle$

 $\rightarrow_k \mathcal{K}\langle E, I_d, W, \{ \mathcal{S}\langle \iota_d, \Sigma'' \rangle, \mathcal{S}\langle \iota_d', \Sigma' \rangle \} \cup S, D \rangle$

(S-GLOBAL)

 $\Sigma = \mathcal{E}[\mathcal{S}_{GLB}(x)] \qquad \Sigma' = \mathcal{E}[E(x)]$ $\mathcal{K}\langle E, I_d, W, \{\mathcal{S}\langle \iota_d, \Sigma\rangle\} \cup S, D\rangle$ $\rightarrow_k \mathcal{K}\langle E, I_d, W, \{ \mathcal{S}\langle \iota_d, \Sigma' \} \cup S, D \rangle$

(S-DEPLOY-NEW)

 $\Sigma = \mathcal{E}[\mathcal{S}_{\mathcal{D}E\mathcal{P}}\langle\iota_b, \mathcal{C}\langle\iota_c, N, \Sigma_c\rangle, \overline{\sigma}\rangle] \qquad (\iota_b, \iota_c) \notin \operatorname{dom}(D) \qquad \iota'_d \operatorname{fresh}$ $w = \mathcal{W}\langle \iota'_d, N, \Sigma_c[in_{i,0} \mapsto \sigma_i \mid \forall i \in [1., |\overline{\sigma}|]] \rangle \qquad \Sigma' = \mathcal{E}[\iota'_d]$ $D' = D[(\iota_b, \iota_c) \mapsto \iota'_d]$ $\mathcal{K}\langle E, I_d, W, \{\mathcal{S}\langle \iota_d, \Sigma\rangle\} \cup S, D\rangle$ $\rightarrow_k \mathcal{K}\langle E, \{\iota'_d\} \cup I_d, \{w\} \cup W, \{\mathcal{S}\langle \iota_d, \Sigma'\rangle\} \cup S, D'\rangle$

(S-DEPLOY-EXISTING) $\Sigma = \mathcal{E}[S_{\mathcal{D}E\mathcal{P}}\langle \iota_b, \mathcal{C}\langle \iota_c, N, \Sigma_c \rangle, \overline{\sigma} \rangle] \qquad \iota'_d = D(\iota_b, \iota_c) \qquad \Sigma' = \mathcal{E}[\iota'_d]$ $\mathcal{K}\langle E, I_d, W, \{\mathcal{S}\langle \iota_d, \Sigma\rangle\} \cup S, D\rangle$ $\rightarrow_k \mathcal{K}\langle E, \{\iota'_d\} \cup I_d, W, \{\mathcal{S}\langle\iota_d, \Sigma'\} \cup S, D\rangle$

Inter-Turn Semantics (Driver Loop)



Karcharias

PLT Redex implementation of Karcharias can be found online:

https://gitlab.soft.vub.ac.be/boeyen/karcharias

R $p \in \mathbf{Program}$::= Local Wiring Rules (\rightarrow_w) : $r \in R \subseteq$ Reactor ::= $\mathcal{R}\langle x,N\rangle$ $(\overline{i}, nt, \overline{o})$ $n \in N \subseteq$ Node ::= (W-REF) $i \in$ Input Port ::= $x \mid v$ $\mathcal{W}\langle \iota_d, \{(\overline{i}_l \& [x] \& \overline{i}_r, nt, \overline{o})\} \cup N, \Sigma \rangle$ Local Snapshot Rules (\rightarrow_s) : $\rightarrow_{w} \mathcal{W}(\iota_{d}, \{(\overline{i}_{l} \in [\Sigma(x)] \in \overline{i}_{r}, nt, \overline{o})\} \cup N, \Sigma)$ $nt \in Node Type$::= RHO(N) (S-SELF-REF) (S-DEPLOY-PRIMITIVE) DEPLOY (W-DEPLOY) $\Sigma = \mathcal{E}[S_{\mathcal{REF}}(\iota_d, x)] \qquad v = \Sigma(x)$ $\begin{array}{c} S\langle \iota_d, \mathcal{E}[S_{\mathcal{D}E\mathcal{P}}\langle \iota_b, p, \overline{v}\rangle] \rangle \\ \rightarrow_s S\langle \iota_d, \mathcal{E}[\delta_p(\overline{v})] \rangle \end{array}$ x, ι_h fresh $S\langle \iota_d, \Sigma \rangle \rightarrow_s S\langle \iota_d, \mathcal{E}[v] \rangle$ $o \in \mathbf{Output Port}$::= x $\Sigma' = \Sigma[x \mapsto \mathcal{S}_{\mathcal{D} \mathcal{E} \mathcal{P}} \langle \iota_b, \sigma, \overline{\sigma} \rangle][o_i \mapsto \mathcal{S}_{\mathcal{R} \mathcal{E} \mathcal{F}} \langle \mathcal{S}_{\mathcal{R} \mathcal{E} \mathcal{F}} \langle \iota_d, x \rangle, out_{i,0} \rangle | \forall i \in [1., |\overline{o}|]]$ (S-TUPLE-REF) $x \in X \subseteq$ Name $\mathcal{W}(\iota_d, \{([\sigma] \ \ \overline{\sigma}, \mathcal{DEPLOY}, \overline{o})\} \cup N, \Sigma) \to_w \mathcal{W}(\iota_d, N, \Sigma')$ $S(\iota_d, \mathcal{E}[S_{\mathcal{REF}}(\overline{v}, out_{(0,i)})])$ Semantic Entities: $\mathcal{K}^{inter}\langle \tau,k\rangle$ $k^{inter} \in$ Inter-Turn Configuration ::= $\rightarrow_s S(\iota_d, \mathcal{E}[v_i])$ $\{in_{i,j}, out_{i,j} \mid \forall i \in \mathbb{N}^+, \forall j \in \mathbb{N}\} \subseteq X$ (W-RHO) $\tau \subseteq$ Primitive Time-Varying Sources ::= $\{x \mapsto \overline{v}, \ldots\}$ Global Rules (\rightarrow_k) : $v \in V \subseteq$ Domain Value $\iota_c \text{ fresh } \Sigma_c = \text{shift}_io(\Sigma) \quad c = \mathcal{C}(\iota_c, N_{inner}, \Sigma_c)$ $\widetilde{k} \in \widetilde{K} \subset K$ where $\widetilde{k} \rightarrow k$ (S-ACTIVATE) $\overline{\mathcal{W}(\iota_d, \{(\overline{\sigma}, \mathcal{RHO}(N_{inner}), [o])\} \cup N, \Sigma\} \to_w \mathcal{W}(\iota_d, N, \Sigma[o \mapsto c])}$ Initial Configuration: (S-CONGRUENCE) $\mathcal{W}\langle \iota_d, \emptyset, \Sigma \rangle \in W \quad \iota_d \in I_d \\ \forall s \in S : s = S \langle \iota'_d, \Sigma' \rangle, \iota_d \neq \iota'_d$ $k_{init}^{inter} = \mathcal{K}^{inter} \langle \tau_{init}, k_{init} \rangle$ $s \rightarrow_s s'$ $in_{i,j+1}$ $x = in_{i,j}$ $\mathcal{K}\langle E, I_d, W, \{s\} \cup S, D \rangle$ $\frac{\mathcal{K}\langle E, I_d, W, S, D \rangle}{\rightarrow_k \mathcal{K}\langle E, I_d, W, \{S(\iota_d, \Sigma)\} \cup S, D \rangle}$ $\mathcal{K}\langle E, I_d, W, S, D \rangle$ Where shift_io(Σ) = { $x' \mapsto \sigma$ | $\forall x \mapsto \sigma \in \Sigma, x' = \{out_{i,j+1} \mid x = out_{i,j}\}$ $k \in K \subseteq$ Configuration $\tau_{init} = \{time \mapsto [0, 1, 2, 3, \ldots], \ldots\}$::= $\rightarrow_k \mathcal{K}(E, I_d, W, \{s'\} \cup S, D)$ otherwise $w \in W \subseteq$ Deployment Wiring ::= $\mathcal{W}(\iota_d, N, \Sigma)$ Reduction Relation (\sim): (S-REF) Global Rules (\rightarrow_{L}) : $S(\iota_d, \Sigma)$ $s \in S \subseteq$ Deployment Snapshot ::= $\Sigma = \mathcal{E}[S_{\mathcal{REF}}(\iota'_d, x)] \qquad v = \Sigma'(x) \qquad \Sigma'' = \mathcal{E}[v]$ (INTRA-TURN) ::= $\mathcal{C}\langle \iota_c, N, \Sigma \rangle$ $\mathcal{K}\langle E, I_d, W, \{ \mathcal{S}\langle \iota_d, \Sigma \rangle, \mathcal{S}\langle \iota_d', \Sigma' \rangle \} \cup S, D \rangle$ $c \in Capture$ (w-congruence) $\rightarrow_k \mathcal{K}\langle E, I_d, W, \{ \mathcal{S}\langle \iota_d, \Sigma'' \rangle, \mathcal{S}\langle \iota_d', \Sigma' \rangle \} \cup S, D \rangle$ $k \rightarrow_k k'$ $\sigma \subseteq \mathbf{Signal}$ $w \rightarrow_w w'$::= 71 $\overline{\mathcal{K}^{inter}(\tau,k) \rightsquigarrow \mathcal{K}^{inter}(\tau,k')}$ (S-GLOBAL) $S_{GLB}\langle x \rangle$ $\mathcal{K}\langle E, I_d, \{w\} \cup W, S, D \rangle \to_k \mathcal{K} \langle E, I_d, \{w'\} \cup W, S, D \rangle$ $\Sigma = \mathcal{E}[\mathcal{S}_{GLB}(x)] \qquad \Sigma' = \mathcal{E}[E(x)]$ $\mathcal{S}_{\mathcal{R}\mathcal{E}\mathcal{F}}(\sigma,x)$ $\mathcal{K}(E, I_d, W, \{S(\iota_d, \Sigma)\} \cup S, D)$ (NEXT-TURN) $\mathcal{S}_{\mathcal{DEP}}(\iota_b,\sigma,\overline{\sigma})$ $\rightarrow_k \mathcal{K} \langle E, I_d, W, \{ \mathcal{S} \langle \iota_d, \Sigma' \rangle \} \cup S, D \rangle$ $k = \mathcal{K} \langle E, I_d, W, S, D \rangle \in \widetilde{K}$ i $\ldots |\sigma|$ $S(\iota_{maind}, E_{main}) \in S$::= (S-DEPLOY-NEW) $\Sigma = \mathcal{E}[S_{DEP}(\iota_b, C(\iota_c, N, \Sigma_c), \overline{\sigma})] \quad (\iota_b, \iota_c) \notin \text{dom}(D) \quad \iota'_d \text{ fresh}$ $k' = \mathcal{K} \langle E[x_i \mapsto v_{i,now} | \forall i \in [1.,|\tau|]], \{\iota_{main.d}\}, W, \emptyset, D \rangle$ 7) ::= $\dots |p| c |\iota_d| \overline{v}$ $w = \mathcal{W}\langle i'_{d}, N, \Sigma_{c}[in_{i,0} \mapsto \sigma_{i} \mid \forall i \in [1..|\overline{\sigma}|]]\rangle \qquad \Sigma' = \mathcal{E}[i'_{d}]$ $\overline{\mathcal{K}^{inter}}\langle\{x_i\mapsto [v_{i,now}]\,\S\,\overline{v}_i\,|\,\forall i\in[1..|\tau|]\},k\rangle$ $D' = D[(\iota_b, \iota_c) \mapsto \iota_d']$ $E \subseteq \Sigma \subseteq$ Value Environment ::= $\{x \mapsto v, \ldots\}$ $\sim \mathcal{K}^{inter} \langle \{ x_i \mapsto \overline{v}_i \mid \forall i \in [1, |\tau|] \}, k' \rangle$ $\mathcal{K}\langle E, I_d, W, \{S\langle \iota_d, \Sigma\rangle\} \cup S, D\rangle$ $\Sigma \subseteq$ Signal Environment ::= $\{x \mapsto \sigma, \ldots\}$ $\rightarrow_k \mathcal{K}\langle E, \{\iota'_d\} \cup I_d, \{w\} \cup W, \{\mathcal{S}\langle\iota_d, \Sigma'\rangle\} \cup S, D'\rangle$ $\{(\iota_h, \iota_c) \mapsto \iota_d, \ldots\}$ $D \subseteq$ Toggle Environment ::= (S-DEPLOY-EXISTING) $\Sigma = \mathcal{E}[\mathcal{S}_{DEP}(\iota_b, \mathcal{C}(\iota_c, N, \Sigma_c), \overline{\sigma})] \quad \iota'_d = D(\iota_b, \iota_c) \quad \Sigma' = \mathcal{E}[\iota'_d]$ $\iota_b \in I_b \subseteq$ Branching Point Id, $\iota_c \in I_c \subseteq$ Capture Id, $\iota_d \in I_d \subseteq$ Deployment Id $\mathcal{K}\langle E, I_d, W, \{S\langle \iota_d, \Sigma \rangle\} \cup S, D \rangle$ $p \in \mathbf{Primitives}, \delta_p : V \star \to V \star \text{(for every } p\text{)}$ $\rightarrow_k \mathcal{K}\langle E, \{\iota'_d\} \cup I_d, W, \{S\langle \iota_d, \Sigma'\} \cup S, D\rangle$ Given $p = \{\mathcal{R}(x_1, N_1), \dots, \mathcal{R}(x_{|p|}, N_{|p|})\}$ (|p| is the total number of user-defined reactor definitions): $k_{init} = \mathcal{K}(E_{init}, \emptyset, W_{init}, \emptyset, \emptyset)$ $\begin{aligned} E_{init} &= \{x_i \mapsto \operatorname{ntc}(N_i) \mid \forall i \in [1..|p|]\} \cup \{time \mapsto 0, \ldots\} \\ W_{init} &= \{\mathcal{W}(\iota_{d,main}, \{([main], DEPLOY, [out_{(i,0)}] \mid \forall i \in [1..|o|_{main}]])\}, \Sigma_{init}\} \} \end{aligned}$ $= \{+ \mapsto [+], \ldots\} \cup \{x_i \mapsto S_{GLB}(x_i) \mid \forall i \in [1, |p|]\} \cup \{time \mapsto S_{GLB}(time), \ldots\}$ Σ_{init} **Inter-Turn Semantics Intra-Turn Semantics** Where $|o|_{main}$ is the number of outputs defined in the main reactor in p, and $\operatorname{ntoc}(N) = \mathcal{C}(\iota_c, N, \Sigma_i nit)$ (where ι_c fresh)

Motivation

Formalisation

Lessons learned

Lessons Learned

- Graph-based RP needs wiring and propagation
 - w-rules and s-rules
 - Wiring and propagating can, in general, be interleaved!
 - Run-time wiring decisions due to higher-order reactivity
- Memory consumption ...

Remember this graph?



Why does this happen?

(>>>): sequential composition of SFs

```
(>>>) :: SF a b -> SF b c -> SF a c
(SF {sfTF = tf1}) >>> (SF {sfTF = tf2}) =
    SF {sfTF = tf}
    where
    tf dt a = (sf1' >>> sf2', c)
    where
      (sf2', c) = tf2 dt b
      (sf1', b) = tf1 dt a
```

Elegant, because it's all *just* composition of functions.

Signal Functions

Actually, this is not what really happens due to optimisations...

Nilsson, H., Courtney, A., & Peterson, J. (2002, October). Functional reactive programming, continued. In *Proceedings of the 2002 ACM SIGPLAN workshop on Haskell* (pp. 51-64).

Without optimisations



Cooper, G. H., & Krishnamurthi, S. (2006, March). Embedding dynamic dataflow in a call-by-value language. In *European symposium on programming* (pp. 294-308). Springer, Berlin, Heidelberg.

GC isn't possible!

Each recursive application of (>>>) extends the function's lexical environment. In addition to Haskell's lazy evaluation model, this introduces a space leak!

> Which is what graph-based RPLs avoid altogether by treating RP programs as graphs [*].

In case of Karcharias...

 $\begin{aligned} & (\text{NEXT-TURN}) \\ & k = \mathcal{K}\langle E, I_d, W, S, D \rangle \in \widetilde{K} \\ & \mathcal{S}\langle \iota_{main,d}, E_{main} \rangle \in S \\ \\ & \frac{k' = \mathcal{K}\langle E[x_i \mapsto v_{i,now} \mid \forall i \in [1..|\tau|]], \{\iota_{main,d}\}, W, \emptyset, D \rangle}{\mathcal{K}^{inter} \langle \{x_i \mapsto [v_{i,now}] \$ \overline{v}_i \mid \forall i \in [1..|\tau|]\}, k \rangle \\ & \sim \mathcal{K}^{inter} \langle \{x_i \mapsto \overline{v}_i \mid \forall i \in [1..|\tau|]\}, k' \rangle \end{aligned}$

- Easy/ier to determine memory consumption
 - Wirings remain static, after they have been constructed.
 - Snapshots are "reduced" versions of wirings (with actual values of signals)

Question now becomes: when do we have to wire?

Future Work

- Currently lacking:
 - No error handling / detection of stuck states
 - Stateful computations aren't formalised yet
 - Thorough comparison between function-based RP and graph-based RP
- Ultimate goal
 - Expressive, higher-order, RP language that works for very small computers
 - Statically classify RP programs w.r.t. their required memory allocation/consumption behaviour



Looking for ideas!

A Graph-Based Formal Semantics of Reactive Programming from First Principles

<u>**Bjarno Oeyen**</u>, Joeri De Koster, Wolfgang De Meuter Software Languages Lab, Vrije Universiteit Brussel, Belgium

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